

# Some additional topics

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October 2, 2019



# Inductive Proofs



## Theorem (Induction)

Let  $P(n)$  be some property depending on  $n \in \mathbb{N}$ . If

1  $P(n_0)$  is true and

2 for every  $n \in \mathbb{N}$   $P(n) \Rightarrow P(n+1)$  is true

then  $P(n)$  holds for all  $n \in \mathbb{N}$ .



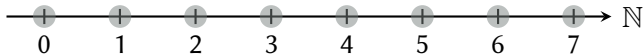
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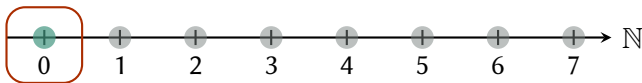
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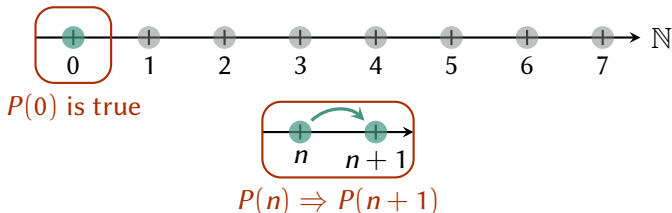
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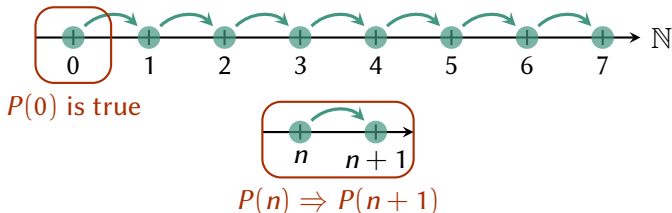
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# Optimization with side constraints

The Karush-Kuhn-Tucker conditions





- With Lagrange Optimization we can solve problems of the form

$$\begin{aligned} \min f(x) \\ \text{s.t. } g_i(x) = 0 \quad \text{for } i = 1, \dots, k \end{aligned}$$

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  - budget constraints
  - time constraints
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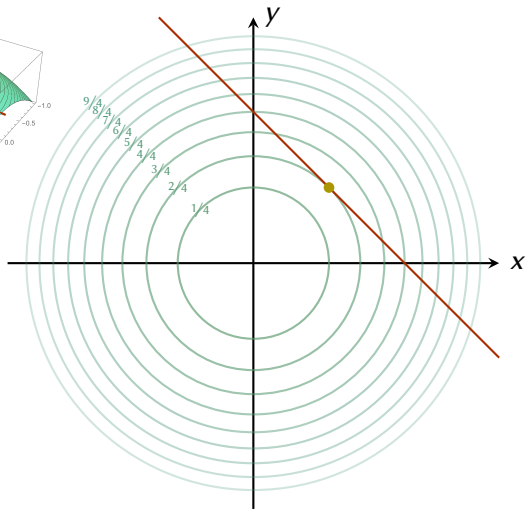
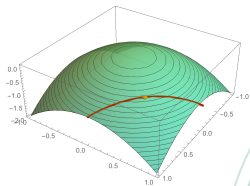
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- Many problems have *inequalities* as side constraints, e.g.
  - budget constraints
  - time constraints
  - capacity constraints
- We want to solve problems of the form

$$\begin{aligned} \min f(x) \\ \text{s.t. } g_i(x) \leq 0 \quad \text{for } i = 1, \dots, k \\ h_j(x) = 0 \quad \text{for } j = 1, \dots, l \end{aligned}$$

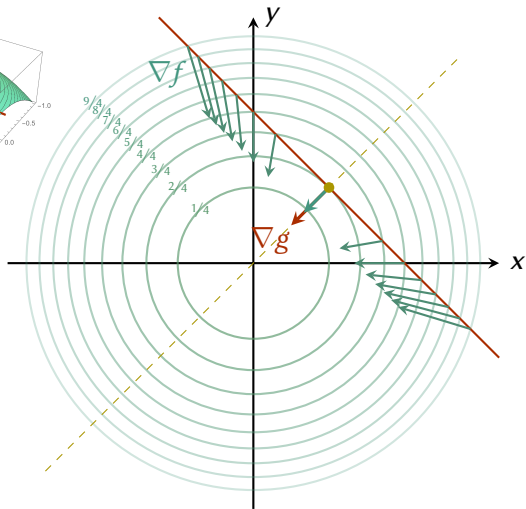
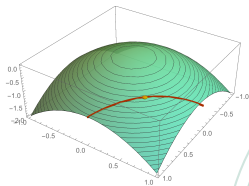


# Lagrange Multipliers



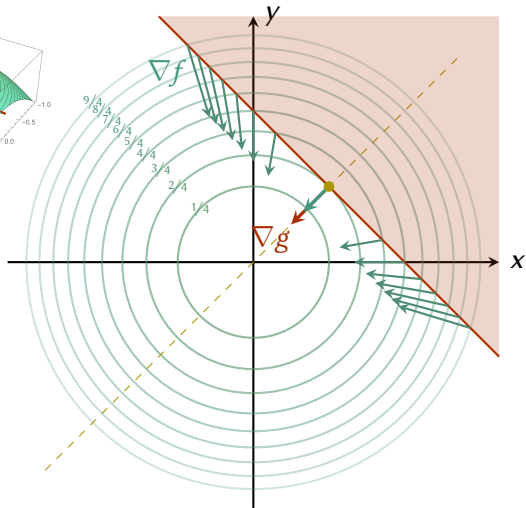
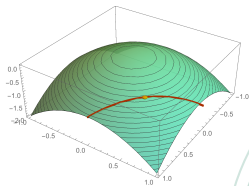
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## Theorem

Let  $\mathbf{x}^{(0)} = (x_n^{(0)}, \dots, x_n^{(0)})$  be an optimal solution of

$$\begin{aligned} \max \quad & f(\mathbf{x}) \\ \text{s.t.} \quad & g_i(\mathbf{x}) \geq 0 \quad \forall i = 1, \dots, k \\ & h_j(\mathbf{x}) = 0 \quad \forall j = 1, \dots, l \end{aligned}$$

and  $f, g, h$  partially differentiable then there are  $\lambda_1, \dots, \lambda_k \in \mathbb{R}$  and  $\mu_1, \dots, \mu_l \in \mathbb{R}$  such that

$$\nabla f(\mathbf{x}^{(0)}) - \sum_{i=1}^k \lambda_i \nabla g_i(\mathbf{x}^{(0)}) - \sum_{j=1}^l \mu_j \nabla h_j(\mathbf{x}^{(0)}) = 0$$

and

$$\begin{aligned} \lambda_j &\geq 0 \quad \forall j = 1, \dots, k \\ \lambda_j h_j(\mathbf{x}^{(0)}) &= 0 \quad \forall j = 1, \dots, k \end{aligned}$$



# Optimization with side constraints

## The Simplex-Algorithm





- An interesting special case of optimization is the following:

$$\begin{aligned} \max f(x) \\ \text{s.t. } g_i(x) \leq 0 \quad \text{for } i = 1, \dots, k \end{aligned}$$

where  $f$  and  $g$  are *linear* functions.



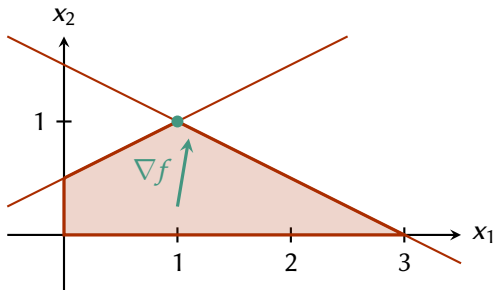
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- Example:

$$\begin{aligned} \max \quad & x_1 + 6x_2 \\ \text{s.t.} \quad & x_1 + 2x_2 \leq 3 \\ & -x_1 + 2x_2 \leq 1 \\ & x_1, x_2 \geq 0 \end{aligned}$$



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$$\begin{aligned} \max \quad & x_1 + 6x_2 \\ \text{s.t.} \quad & x_1 + 2x_2 + y_1 = 3 \\ & -x_1 + 2x_2 + y_2 = 1 \\ & x_1, x_2, y_1, y_2 \geq 0 \end{aligned}$$

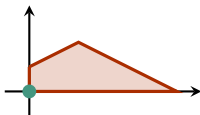


# The simplex method

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0	-1	-6	0	0
3	1	2	1	0
1	-1	2	0	1



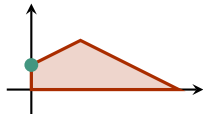
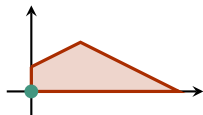
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$$\begin{array}{c|cccc} 0 & -1 & -6 & 0 & 0 \\ \hline 3 & 1 & 2 & 1 & 0 \\ 1 & -1 & 2 & 0 & 1 \end{array}$$

$$\begin{array}{c|cccc} 3 & -4 & 0 & 0 & 3 \\ \hline 2 & 2 & 0 & 1 & -1 \\ 1/2 & -1/2 & 1 & 0 & 1/2 \end{array}$$

$$\begin{aligned} \max \quad & x_1 + 6x_2 \\ \text{s.t.} \quad & x_1 + 2x_2 + y_1 = 3 \\ & -x_1 + 2x_2 + y_2 = 1 \\ & x_1, x_2, y_1, y_2 \geq 0 \end{aligned}$$



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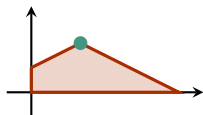
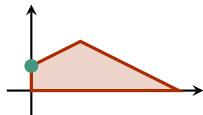
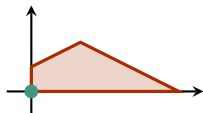
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3	1	2	1	0
1	-1	2	0	1

3	-4	0	0	3
2	2	0	1	-1
1/2	-1/2	1	0	1/2

7	0	0	2	1
1	1	0	1/2	-1/2
1	0	1	1/4	1/4



# Conclusion





Things you...	... should know	... should keep in mind
Linear Algebra	Basic vector and matrix computations Solving systems of linear equations	Eigenvalues
Analysis	Derivatives Basic integration	Induction Sequences & Limits
Optimization	Computing extrema of uni-/multivariate functions Lagrange Optimization	KKT conditions Simplex Algorithm



# Thank you!

And good luck with your studies at HU Berlin!

